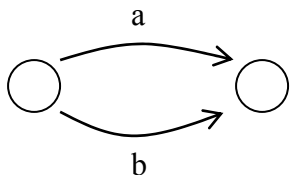


# Regular Expressions and Finite Automata

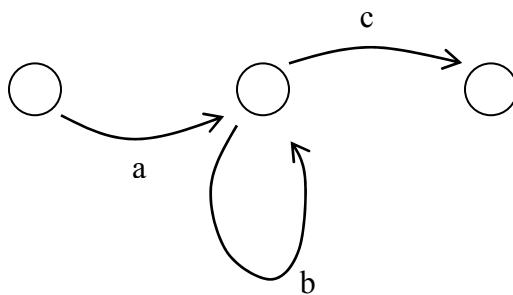
This text shows how to convert a regular expression into a graphical representation of a finite automata and then to an algorithm that analyzes this automata. However, the techniques shown are informal — some regular expressions cannot be transformed by them.

Each of the items below show how to convert a regular expression into the graphical representation of a finite automata.

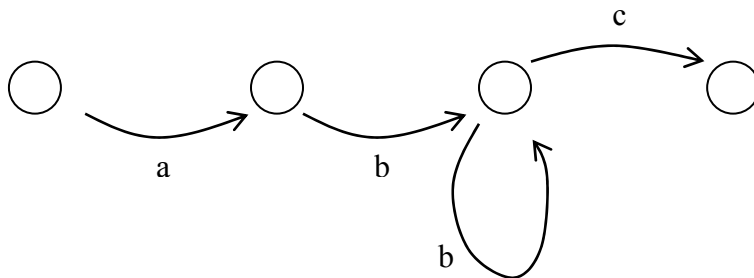
a) or. Example:  $a \mid b$



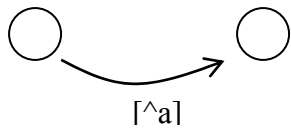
b) zero or more repetitions. Example:  $ab^*c$



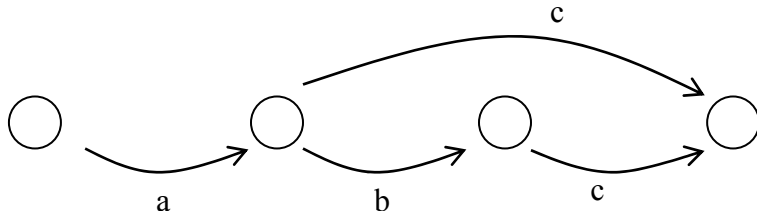
c) One or more repetitions. Example:  $ab^+c$



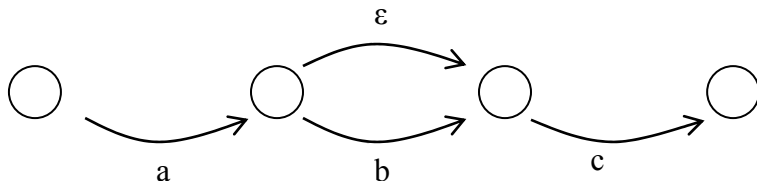
d) Everything but something. Example:  $[\wedge a]$



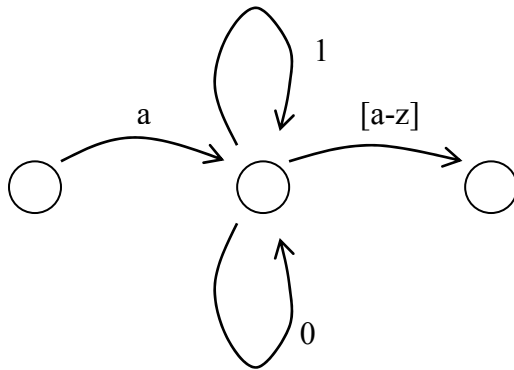
e) Option. Example:  $ab?c$



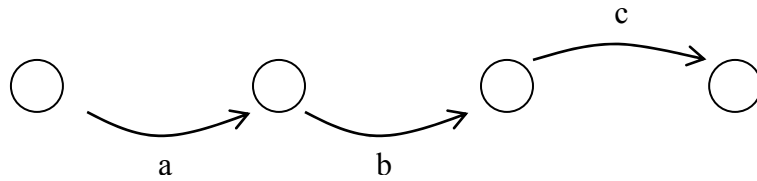
or



f) Repetition of optional things. Example:  $a(0|1)^*[a-z]$



g) Sequence of things. Example:  $abc$



Now we show how to convert a regular expression into a Java method that recognizes the expression. The example should explain itself.

An example of a regular expression would be:

`a(0|1)*#[a-z](a|b)0*`

The analyzer to recognize it would be:

```

void analyze(char in[]) {
    int s = 1;
    int k = 0;
    while ( s != 6 ) {
        switch (s) {
            case 1:
                if ( in[k] != 'a' )
                    error();
                else {
                    k++;
                    s = 2;
                }
                break;
            case 2:
                if ( in[k] == '0' || in[k] == '1' )
                    k++;
                else if ( in[k] == '#' ) {
                    k++;
                    s = 3;
                }
                else
                    error();
                break;
            case 3:
                if ( in[k] == '#' )
                    k++;
                else if ( in[k] >= 'a' && in[k] <= 'z' ) {
                    k++;
                }
        }
    }
}
  
```

```

        s = 4;
    }
    else
        error();
    break;
case 4:
    if ( in[k] == 'a' || in[k] == 'b' ) {
        k++;
        s = 5;
    }
    else
        error();
    break;
case 5:
    if ( in[k] == '0' )
        k++;
    else
        s = 6;
    break;
}
}
}

```

Caveat: some regular expressions cannot be converted into a finite automaton by the techniques given in this manual. For example, the method fails when using the regular expressions

```

b*b
a+a
(a|b)+ab
b?b

```

Why? In  $b^*b$ , the automaton would advance through all  $b$ 's in the input, consuming even the  $b$  that would match the last  $b$  in " $b^*b$ ".